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Flexible Schema Data (FSD) Management in RDBMS Opportunities & Challenges for NoSQL

Zhen Hua Liu Architect
Dieter Gawlick Architect

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Program Agenda

- 1** ➤ **DBMS Generational Gap: Need Support of Flexible Schema Data**
- 2** ➤ Requirements & Challenges of FSD management
- 3** ➤ Engineering Practices & Principles of FSD management
- 4** ➤ Each FSD Principle with Opportunities & Challenges Analysis
- 5** ➤ Related Work & Conclusion
- 6** ➤ Q & A

DBMS Generational Gap

Upfront Schema Design becomes bottleneck!

- My Father's RDBMS: Classical Bank Record Keep Application
 - Process Statically Shaped Data
 - DBA Style: Schema First, Data Later
 - Every New Shaped Data demands Schema Evolution
- My Generation's DBMS: Flexible Schema Data Management
 - Process Semi-Structured/Unstructured Dynamically Shaped Data: Web Data, Diversity Data
 - Agile Style: Data First, Schema Later/Never
 - **On Write: No Schema for Store; On Read: Soft Schema for Query**

Simple & Popular FSD Example - JSON

```
{
  "firstName": "John",
  "lastName": "Smith",
  "age": 25,
  "address": {
    "streetAddress": "21 2nd Street",
    "city": "New York",
    "state": "NY",
    "postalCode": "10021",
    "isBusiness": false
  },
  "phoneNumbers": [
    {"type": "home", "number": "212 555-1234"},
    {"type": "fax", "number": "646 555-4567"}
  ],
  "creditHistory": [
    {"year": 2011, "creditScore": 650}
  ],
  "bankruptcies": null
}
```

- Each FSD instance is a set of key-value pairs organized hierarchically
- FSD collection: a set of FSD instances
- Schema is considered **unbounded** for the entire FSD collection
- Storing FSD using relational tables requires **Constant Schema Evolution**
- Storing FSD using vertical table requires many self-joins & suboptimal object retrieval time

FSD Management Requirements

Conventional Schema-based RDBMS Wisdoms are being Challenged!

- **Storage Requirement:** How to store data without upfront schema definition ?
- **Query Requirement:** How to query data without upfront schema definition ?
- **Indexing Requirement:** How to index data without upfront schema definition ?

NoSQL ? Does this imply SQL is dead ?

- **Consolidated Data Management Platform Requirement:** How to query both my data and my father's data together ?

Think Out of Box here means **THINK OUT OF SCHEMA !**

New Engineering Practises/Principles for managing FSD

My Data & My Father's Data are jointly queried together via NoSQL

- **Storage Principle: Schema-less Storage**
 - **Document Object Storage Model:** Not relying on schema to decompose & shred data.
 - **DataGuide:** dynamically computed SOFT Schema to support **schema on read** capability
- **Query Principle: Declarative Query (Central Dogma of DBMS)**
 - **NoSQL:** Naturally open SQL as **Set Query Language** with embedded FSD domain language
 - **Declarative FSD Domain Language:** for query both FSD schema & data together
- **Index Principle: Schema-less Indexing**
 - **Search Index :** Data First/Schema Never Indexing – Ad-hoc workload
 - **Table index:** Data First/Schema Later Indexing – known workload

My Principles are inspired from My Father's RDBMS Extensibility Ideas (UDT/UDF/UDI)

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FSD Storage principle in RDBMS

Storage Principle – Document Object Store without shredding

- Each FSD instance is **self-contained** without relying on central schema definition
 - Schema & Data are stored together in each FSD instance
- Each domain specific FSD can be stored as varchar/varbinary/CLOB/BLOB **without a new SQL datatype!**
 - JSON as FSD Example: FSD Check Constraint/Soft Schema Validation
 - *CREATE TABLE PERSON_JTAB (jcol VARCHAR(32000) CONSTRAINT jcon CHECK (jcol IS JSON));*
- 100% operational completeness support for FSD
 - Transactions, Replication, Partition, Security, Temporal, Provenance, Export/Export, Fault Tolerance, Client APIs, ...

FSD Query Principle in RDBMS

Query Principle: SQL: Set Query Language with FSD Domain Language

- Leveraging SQL as **Inter-Document Set Query Language**
- Leveraging FSD Domain Language as **Intra-Document** query language
- FSD Domain language for each FSD instance to do
 - path navigation, extracting scalar values & fragments, searching content, transforming and updating fragments. JSON as an example

```
SELECT JSON_VALUE(T.jcol, '$.person.name'),  
       JSON_QUERY(T.jcol, '$.person.address')  
FROM PERSON_JTAB T  
WHERE JSON_EXISTS(T.jcol, '$.person.creditHistory?(score >=  
700)') AND  
JSON_TEXTCONTAINS(T.jcol, '$.person..experiences', 'semi-  
structured data processing')  
ORDER BY JSON_VALUE(T.jcol, '$.person.address.zip')
```

FSD Query Principle in RDBMS

JSON_TABLE() for Relational View Projection – UNNEST array

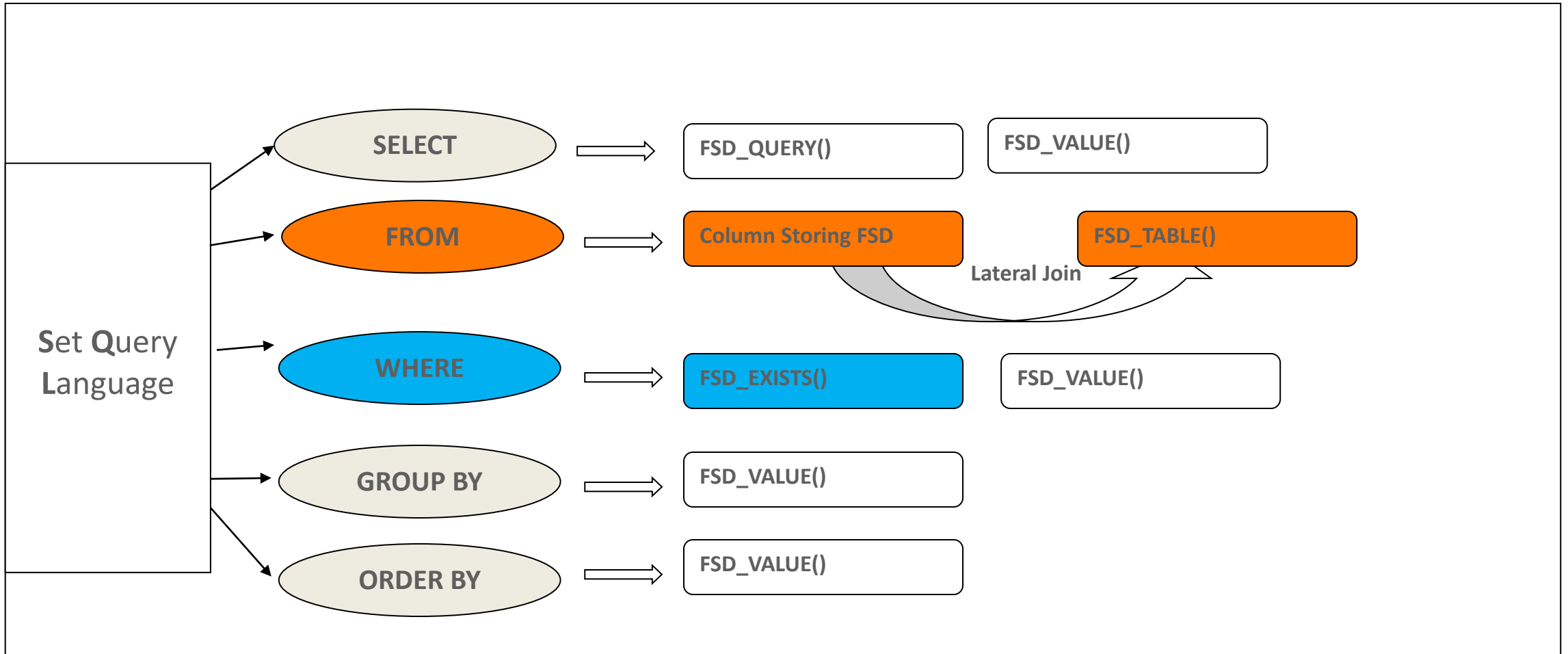
- Project FSD data as relational table
- Objects inside array become ROWS
- Values inside object become COLUMNS

PH_TYPE	PH_NUM
home	212 555-1234
fax	646 555-4567

```
SELECT jrv.*
FROM Person_Jtab, JSON_TABLE (jcol, '$.person.phoneNumbers[*]'
                             COLUMNS
                             PH_TYPE    VARCHAR2(10) PATH '$.type' ,
                             PH_NUM     VARCHAR2(10) PATH '$.number') jrv;
```

FSD Query Principle in RDBMS

FSD UDFs to Set Query Language without a brand new query language!



FSD Index principle in RDBMS

Search Index: Data First/Schema Never Indexing

- Schema and data are indexed together for ad-hoc exploratory and search queries over FSD collection
- Generalized Inverted Index (FSD_EXISTS(), FSD_TEXTCONTAINS())
 - Classical SIGIR: Full Text Search for Document Content
 - **Beyond classical SIGIR:**
 - Range Value Search for leaf Scalar Value (auto detecting and indexing number, dates, timestamps)
 - Hierarchical Path Containment for both full text search & range value

```
SELECT JSON_VALUE(T.jcol, '$.person.name')
FROM PERSON_JTAB T
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structured data processing')
```

FSD Index principle in RDBMS

Table Index: Data First/Schema Later as Indexing

- Table Index based on Dataguide & known query workload
- FSD_VALUE() Scalar Value Functional Index
- FSD_TABLE() Table Index/Materialized view
- Table Index brings relational model back into FSD as secondary view & indexing structures instead of primary storage structures
 - Provide maximum flexibility because table index is secondary structure that can be dropped and recreated without impacting primary storage structures
 - No schema evolution & management Issue
 - Enables Paradigm of **Data First, Schema Later as table Index**

FSD Data Model Challenges

What's wrong with a single tree model ?

- **Single Hierarchy issue** for document storage Model
 - Document Storage Model imposes single hierarchy restriction whereas relational Model provides flexible hierarchies access
- Your father can do multi-hierarchical access of relational data & you are re-inventing IMS !
- Need DataGuide to E/R Model EcoSystem ?
 - Given students taking courses hierarchy, it shall infer courses being taken by students hierarchy
- Need Declarative Hierarchy Transformation Language for FSD
 - Leverage Category Theory Providing Hierarchy Equivalency Transformation & Query Access ?

FSD Layout Challenges

Row & Columnar Dual FSD Data Format

- **FSD Instance layout** for both efficient path query and piece-wise update (Avro, BSON, Tree encoding)
- **FSD Set Columnar layout of FSD** for Efficient Vector based Set processing query over FSD collection (Parquet, Dremel)
- Which is Better Storage Model ?
- Leveraging idea of InSitu Query Processing for FSD
 - **Exploit instance/set query friendly Data Layout as secondary Just-In-Time In-Memory structures!**
 - **Never Getting Stale with a storage model**

FSD Layout Challenges

Keeping Dual Formats In-Synch in Real Time

- Tension between dual formats:
 - Ingestion/update friendly instance oriented format
 - Optimal for OLTP workload
 - Not so good for query
 - Search & Analytic query friendly set oriented format:
 - columnar & inverted index favors compression & Batch loading, good for OLAP query
 - Not optimal for OLTP workload
- Keeping dual formats transactionally consistent – LSF/LSM/MVCC

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Related Work

Common Theme: Break strong dependency of using schema to store data

- Research:
 - Stanford LOREL Semi-Structured DB
 - XMLDB
 - Argo/SQL, NoBench benchmark from University of Wisconsin
 - SiNew/TeraData from Yale University Research
- Industry:
 - DB2, Oracle, Microsoft, TeraData SQL/JSON, SQL/XML Support
 - No-SQL DB Products: MongoDB, MarkLogic

Conclusion

Take Home Message for FSD: What would happen if I do not depend on schema to store data ?

- Engineering Practises & Principles for FSD
 - **Storage** – Native Store without shredding, Just-in-time secondary structures for both instance & set friendly access patterns
 - **Query** – SQL As Set Query Language with FSD domain Language embedded
 - **Index** – Search Index and Table index
- The underlying philosophy is very simple:
 - **Treat Schema as if it were data:** Store, Index and Query schema along with the data
 - ***My Data & My Father's Data can be queried together by Naturally opening Set Query Language (NoSQL) !***



Questions



Hardware and Software Engineered to Work Together